```
.text
main:
       li $s1, 1
                            # 1
outer:
       li $s2, 1
                           # 2
                                   ##### begin outer loop
inner:
       li $v0, 1
                           # 3
                                         ###### begin inner loop
       move $a0, $s2
                           # 4
       syscall
                           # 5
                                   #
                                         #
       addi $s2, $s2, 1
                           # 6
       slti $t0, $s2, 11
                           # 7
       bnez $t0, inner
                           # 8
                                  #
                                         ###### end of inner loop
       addi $s1, $s1, 1
                           # 9
       slti $t0, $s1, 11
                           # 10
       bnez $t0, outer
                           # 11
                                ##### end of outer loop
       li $v0, 10
                           # 12
       syscall
                           # 13
# Behavior of instruction #8, which is a conditional branch instruction:
       taken 9 times
       fall thru once
      taken 9 times
#
      fall thru once
                        There are 10 outer iterations, so the branch
#
                         is taken 90 times and falls thru 10 times.
      taken 9 times
       fall thru once
# If our branch prediction strategy is "always assume fall thru", then we are
# correct 10 times out of 100 (10%), which is pretty poor performance.
# If our branch prediction strategy is to use a 1-bit predictor:
# We initially assume the branch is fall thru.
#
  #
#
  wrong: *
# So we will be wrong 20 times == 80/100 correct = 80%.
# If our strategy uses a 2-bit predictor, again assume initially fall thru.
#
  #
#
  wrong:
# So we will be wrong 12 times == 88/100 correct = 88%.
```

loop.s -- example for branch prediction in hardware.